### Introduction

Main challenge in maximizing processor utilization for distributed computing is to reduce idling due to synchronization across processors. Synchronization is necessary after every iteration, however, it prevents many numerical algorithms from scaling with number of processors.

We relax this requirement by synchronizing at a lower rate, process referred to as Lazy Synchronization. We present a novel approach to solve parallel Quadratic Programming problems using the proposed *Lazily Synchronized* Dual Ascent (LSDA) algorithm. We also provide optimal rate for synchronization which ensures faster convergence to the solution compared to *Tightly Synchronized Dual Ascent* (TSDA) algorithm.

### **Problem Formulation**

$$\min_{x} \sum_{i=1}^{N} f_i(x_i)$$
  
subject to  $\sum_{i=1}^{N} A_i x_i = b$ 

where,  $x_i \in \mathbb{R}^{n/N}$ ,  $A_i \in \mathbb{R}^{m \times \frac{n}{N}}$ ,  $b \in \mathbb{R}^m$ , N: number of cluster nodes,

**QP** Problem:

$$f_i(x_i) = \frac{1}{2} x_i^T Q_i x_i + c_i^T x_i$$
  
where,  $c_i \in \mathbb{R}^{n/N}$  and  $Q_i \in \mathbb{R}^{\frac{n}{N} \times \frac{n}{N}}$  is a SPD matrix

#### LSDA Method:

Step 1: For each cluster node i=1...N,

$$x_{i}^{k+1} = \underset{x_{i}}{\operatorname{argmin}} L(x_{i}, y^{k}) = -Q_{i}^{-1}(A_{i}^{T}y^{k} + c_{i})$$

Step 2: At master node,

$$y^{k+1} = y^k + \eta \sum_{i=1}^{N} (A_i x_i^{tP+1} - \frac{b}{N})$$

where, L() is Langrangian,  $\eta$  is step size, y is dual variable, k is iteration such that,  $tP \le k < (t+1)P$ ,  $t \in$  and  $P \ge 1$  is synchronization period.

Optimal Synchronization Period  $(P^*)$ :

$$P^* = \max \arg \min_{P \in \mathbb{N}} \max\{|1 - \lambda_{\min}(M)P|, |1 - \lambda_{\max}(M)P|\}$$

where,  $M = \eta \sum_{i=1}^{N} A_i Q_i^{-1} A_i^T$ ,  $\lambda_{\min}(M)$  and  $\lambda_{\max}(M)$  denote the minimum and the maximum eigenvalues of the square matrix M, respectively.

# A Relaxed Synchronization Approach for Solving Parallel Quadratic Programming Problems with Guaranteed Convergence

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## LSDA Framework



## Results

#### **1. Optimal Synchronization Period**, *P*\*



For a given synthetic dataset, minimum number of iterations k required for LSDA to converge occurs at  $P^* = 70$ , which is independent of the cluster size N. It also validates the analytical formulation for optimal synchronization period

#### 2. Convergence of LSDA vs TSDA

It is observed that LSDA  $(P^* = 70)$ converges to the optimal solution in iterations, much faster *k*=211 compared to k=868 iterations for TSDA. In LSDA, synchronization among nodes occurs only thrice  $(k/P^*)$ whereas in TSDA it occurs every iteration. Hence, the communication delay is reduced by 99.65% in LSDA.



#### **3. Overall Execution time and Speedup**

Computation time and communication time (synchronization), together is defined as the overall execution time. In Figures below, (a) LSDA technique provides significant benefits in both the computation and communication time compared to (b) TSDA method. This is because computationally intensive local x-update calculation and communication demanding synchronization occurs only at optimal synchronization period  $P^*=70$  for LSDA. TSDA performs slower since its synchronization period  $P^*=1$ . (c) Speedup of LSDA vs TSDA.

## Current Research

**QRSVM:** QR decomposition framework to solve large-scale linear *Support* Vector Classification problems having huge number of instances with relatively smaller dimensionality, i.e.  $d \ll n$ .



QRSVM framework comprises of two main stages, namely, 1) *QR decomposition* of the original input matrix X<sup>^</sup> into Householder reflectors and a matrix R, and 2) *Dual Ascent method* to solve the QRSVM problem for obtaining the normal w to the hyperplane and identifying set of support vectors.

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